Effective results for unit equations over finitely generated domains

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History

Let A be a *finitely generated domain over* $\mathbb{Z}$, that is a commutative integral domain containing $\mathbb{Z}$ which is finitely generated as a $\mathbb{Z}$-algebra.

We have $A = \mathbb{Z}[z_1, \ldots, z_r] \supset \mathbb{Z}$ with the z_i algebraic or transcendental over $\mathbb{Q}$.

Denote by A^* the unit group of A.

Theorem (Siegel, Mahler, Parry, Lang)

Let a, b, c be non-zero elements of A. Then the equation

(1) $$ax + by = c \quad \text{in } x, y \in A^*$$

has only finitely many solutions.

Siegel (1921): $A = O_K$ =ring of integers in a number field K,
Mahler (1933): $A = \mathbb{Z}[1/p_1 \cdots p_t]$, p_i primes,
Parry (1950): $A = O_S$ =ring of S-integers in a number field K,
Lang (1960): A arbitrary finitely generated domain over $\mathbb{Z}$

The proofs of Siegel, Mahler, Parry, Lang are *ineffective*.

Thue equations

Let $A = \mathbb{Z}[z_1, \ldots, z_r] \supset \mathbb{Z}$ be a finitely generated domain over $\mathbb{Z}$, and K its quotient field.

Theorem

Let $F(X, Y) = a_0X^n + a_1X^{n-1}Y + \cdots + a_nY^n \in A[X, Y]$ be a square-free binary form of degree $n \geq 3$ and $\delta \in A \setminus \{0\}$. Then

$$F(x, y) = \delta \quad \text{in } x, y \in A \tag{2}$$

has only finitely many solutions.

Thue equations

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(2) $$F(x, y) = \delta \quad \textit{in } x, y \in A$$

has only finitely many solutions.

Idea of proof.

Assume wlog $a_0 \neq 0$ and factor F in a finite extension of K as $F = a_0 \prod_{i=1}^{n}(X - \beta_i Y)$. Take $B = A[a_0^{-1}, \delta^{-1}, \beta_1, \ldots, \beta_n]$. Then for any solution (x, y) of (2) we have

$$(\beta_2 - \beta_3)\frac{x - \beta_1 y}{x - \beta_3 y} + (\beta_3 - \beta_1)\frac{x - \beta_2 y}{x - \beta_3 y} = \beta_2 - \beta_1, \quad \frac{x - \beta_1 y}{x - \beta_3 y}, \frac{x - \beta_2 y}{x - \beta_3 y} \in B^*.$$

□

Effective results for S-unit equations (I)

Let K be an algebraic number field and $S = \{\mathfrak{p}_1, \ldots, \mathfrak{p}_t\}$ a finite set of prime ideals of O_K. Define $O_S = O_K[(\mathfrak{p}_1 \cdots \mathfrak{p}_t)^{-1}]$.
Then O_S^* consists of all elements of K composed of prime ideals from S.

For $\alpha \in \overline{\mathbb{Q}}$ with minimal polynomial $a_0 X^d + \cdots + a_d \in \mathbb{Z}[X]$ with $\gcd(a_0, \ldots, a_d) = 1$, we define its logar. height $h(\alpha) := \log \max_i |a_i|$.

Theorem (Győry, 1979)

Let $a, b, c \in O_S \setminus \{0\}$. There is an effectively computable number C depending on K, S, a, b, c, such that for every pair x, y with

$$ax + by = c, \quad x, y \in O_S^* \tag{3}$$

we have $h(x), h(y) \leq C$.

Thus, given (suitable representations for) K, S, a, b, c, one can determine effectively (suitable representations for) the solutions of (3).

Proof.

Lower bounds for linear forms in ordinary and p-adic logarithms (Baker, Coates, van der Poorten, Yu). □

Effective results for S-unit equations (II)

Let K be an algebraic number field, $S = \{\mathfrak{p}_1, \dots, \mathfrak{p}_t\}$ a finite set of prime ideals of O_K, and $a, b, c \in O_S \setminus \{0\}$.

Suppose that $[K : \mathbb{Q}] = \delta$, K has discriminant Δ, $\max_i N_{K/\mathbb{Q}}\mathfrak{p}_i \le P$, and $\max\big(h(a), h(b), h(c)\big) \le h$.

Theorem (Győry, Yu, 2006; weaker version)

For every pair x, y with

$$ax + by = c, \quad x, y \in O_S^*$$

we have $h(x), h(y) \le C$ with

$$C = 2^{35}(\delta(\delta + t))^{2(\delta+t)+5}|\Delta|^{1/2}(\log|2\Delta|)^{\delta}\, P^{t+1}(h+1).$$

Unit equations over arbitrary finitely generated domains

In 1983/84 Győry extended his effective result on S-unit equations from 1979 to an effective result for equations

$$ax + by = c \quad \text{in } x, y \in A^*$$

for a special class of finitely generated domains $A = \mathbb{Z}[z_1, \ldots, z_r]$ with some of the z_i transcendental.

Aim:
Prove an effective result for unit equations over *arbitrary* finitely generated domains over $\mathbb{Z}$.

The general effective result

Let $A = \mathbb{Z}[z_1, \ldots, z_r] \supset \mathbb{Z}$ be an arbitrary finitely generated domain over $\mathbb{Z}$. The ideal

$$I := \{f \in \mathbb{Z}[X_1, \ldots, X_r] : f(z_1, \ldots, z_r) = 0\}$$

is finitely generated, say $I = (f_1, \ldots, f_m)$. Thus,

$$A \cong \mathbb{Z}[X_1, \ldots, X_r]/(f_1, \ldots, f_m).$$

Remark. A domain, $A \supset \mathbb{Z} \iff$
$f_1, \ldots, f_m$ generate a prime ideal of $\mathbb{Q}[X_1, \ldots, X_r]$ not containing 1.
There are various algorithms to check this for given $f_1, \ldots, f_m$.

By a *representative* for $a \in A$, we mean a polynomial $f \in \mathbb{Z}[X_1, \ldots, X_r]$ such that $a = f(z_1, \ldots, z_r)$.

Theorem 1 (Győry, E., to appear)

Given $f_1, \ldots, f_m$ and representatives for a, b, c, one can effectively determine representatives for all solutions of

$$ax + by = c \quad \text{in } x, y \in A^*.$$

A quantitative result

Let $A \cong \mathbb{Z}[X_1, \dots, X_r]/(f_1, \dots, f_m)$ be a domain with $A \supset \mathbb{Z}$ and $a, b, c \in A \setminus \{0\}$. For $f \in \mathbb{Z}[X_1, \dots, X_r]$ define

$$\begin{aligned} &\deg f := \text{total degree of } f, \\ &h(f) := \log \max |\text{coefficients of } f| \quad \text{(logarithmic height)}, \\ &s(f) := \max(1, \deg f, h(f)) \quad \text{(size)}. \end{aligned}$$

Theorem 2 (Győry , E.)

Let $\widetilde{a}, \widetilde{b}, \widetilde{c} \in \mathbb{Z}[X_1, \dots, X_r]$ be representatives for a, b, c. Suppose that $f_1, \dots, f_m$, $\widetilde{a}, \widetilde{b}, \widetilde{c}$ have total degrees at most d and logarithmic heights at most h. Then each solution x, y of

$$ax + by = c \quad \text{in } x, y \in A^*$$

has representatives $\widetilde{x}, \widetilde{y}$ such that

$$s(\widetilde{x}), s(\widetilde{y}) \leq \exp\left\{(d+2)^{\kappa^r}(h+1)\right\},$$

where κ is an effectively computable absolute constant > 1.

Theorem 2 $\Longrightarrow$ Theorem 1 (I)

We need the following result:

Theorem (Aschenbrenner, 2004)

Let $f_1, \dots, f_m, b \in \mathbb{Z}[X_1, \dots, X_r] \setminus \{0\}$ of total degrees at most d and logarithmic heights at most h. Suppose there are $g_1, \dots, g_m$ such that

$$(4) \qquad g_1 f_1 + \cdots + g_m f_m = b, \quad g_1, \dots, g_m \in \mathbb{Z}[X_1, \dots, X_r].$$

Then there are such $g_1, \dots, g_m$ with

$$\left.\begin{array}{rcl} \deg g_i & \leq & (d+2)^{\kappa^{r \log(r+1)}}(h+1), \\ h(g_i) & \leq & (d+2)^{\kappa^{r \log(r+1)}}(h+1)^{r+1} \end{array}\right\} \text{ for } i = 1, \dots, m$$

where κ is an effectively computable absolute constant > 1.
Hence it can be decided effectively whether (4) is solvable.

This is an analogue of earlier results of Hermann (1926) and Seidenberg (1972) on linear equations over $F[X_1, \dots, X_r]$, F any field.

Theorem 2 $\Longrightarrow$ Theorem 1 (II)

Corollary (Ideal membership algorithm)

Given $f_1, \ldots, f_m, b \in \mathbb{Z}[X_1, \ldots, X_r]$ *it can be decided effectively whether* $b \in (f_1, \ldots, f_m)$.

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Corollary (Ideal membership algorithm)

Given $f_1, \ldots, f_m, b \in \mathbb{Z}[X_1, \ldots, X_r]$ *it can be decided effectively whether* $b \in (f_1, \ldots, f_m)$.

Corollary (Unit decision algorithm)

Given $b, f_1, \ldots, f_m \in \mathbb{Z}[X_1, \ldots, X_r]$ *it can be decided effectively whether* b *represents a unit of* $A = \mathbb{Z}[X_1, \ldots, X_r]/(f_1, \ldots, f_m)$.

Theorem 2 $\Longrightarrow$ Theorem 1 (II)

Corollary (Ideal membership algorithm)

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Proof.

b represents a unit of A

$\Longleftrightarrow$

there is $b' \in \mathbb{Z}[X_1, \ldots, X_r]$ such that $b \cdot b' \equiv 1 \pmod{(f_1, \ldots, f_m)}$

$\Longleftrightarrow$

there are $b', g_1, \ldots, g_m \in \mathbb{Z}[X_1, \ldots, X_r]$ with $b' \cdot b + g_1 f_1 + \cdots + g_m f_m = 1$. □

Theorem 2 $\Longrightarrow$ Theorem 1 (III)

Let $f_1, , \ldots, f_m$ such that $A \cong \mathbb{Z}[X_1, \ldots, X_r]/(f_1, \ldots, f_m)$, and let $\widetilde{a}, \widetilde{b}, \widetilde{c}$ be representatives for $a, b, c \in A$.

By Theorem 2 there is an effectively computable C such that each solution x, y of

$$(1) \qquad ax + by = c, \quad x, y \in A^*$$

has representatives $\widetilde{x}, \widetilde{y}$ of size $\leq C$.

Theorem 2 $\Longrightarrow$ Theorem 1 (III)

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By Theorem 2 there is an effectively computable C such that each solution x, y of

$$ax + by = c, \quad x, y \in A^* \tag{1}$$

has representatives $\widetilde{x}, \widetilde{y}$ of size $\leq C$.

One can find a representative for each solution of (1) as follows:

Check for each pair $\widetilde{x}, \widetilde{y} \in \mathbb{Z}[X_1, \ldots, X_r]$ of size $\leq C$ whether

$$\widetilde{a} \cdot \widetilde{x} + \widetilde{b} \cdot \widetilde{y} - \widetilde{c} \in (f_1, \ldots, f_m),$$

$$\widetilde{x}, \widetilde{y} \text{ represent elements of } A^*.$$

From the pairs $(\widetilde{x}, \widetilde{y})$ satisfying this test, select a maximal subset of pairs that are different modulo $(f_1, \ldots, f_m)$. $\square$

Exponential equations

Let $A \cong \mathbb{Z}[X_1, \ldots, X_r]/(f_1, \ldots, f_m)$ be a domain with $A \supset \mathbb{Z}$, $a, b, c \in A \setminus \{0\}$, and $\gamma_1, \ldots, \gamma_s$ *multiplicatively independent* elements of $A \setminus \{0\}$. Consider

$$(5) \qquad a\gamma_1^{u_1} \cdots \gamma_s^{u_s} + b\gamma_1^{v_1} \cdots \gamma_s^{v_s} = c \quad \text{in } u_1, \ldots, v_s \in \mathbb{Z}.$$

Theorem 3 (Győry, E.)

Let $\widetilde{a}, \widetilde{b}, \widetilde{c}$, $\widetilde{\gamma}_1, \ldots, \widetilde{\gamma}_s \in \mathbb{Z}[X_1, \ldots, X_r]$ be representatives for $a, b, c, \gamma_1, \ldots, \gamma_s$ and assume that $f_1, \ldots, f_m$, $\widetilde{a}, \widetilde{b}, \widetilde{c}$, $\widetilde{\gamma}_1, \ldots, \widetilde{\gamma}_s$ have total degrees at most d and logarithmic heights at most h. Then for each solution of (5) we have

$$\max(|u_1|, \ldots, |v_s|) \leq \exp\left\{(d+2)^{\kappa^{r+s}}(h+1)\right\}$$

where κ is an effectively computable absolute constant > 1.

An effective criterion for multiplicative (in)dependence

Let $f_1, \ldots, f_m$ be such that $A \cong \mathbb{Z}[X_1, \ldots, X_r]/(f_1, \ldots, f_m)$, let $\gamma_1, \ldots, \gamma_s$ be non-zero elements of A, and choose representatives $\tilde{\gamma}_1, \ldots, \tilde{\gamma}_s$ for $\gamma_1, \ldots, \gamma_s$.

Suppose that $f_1, \ldots, f_m$, $\widetilde{\gamma}_1, \ldots, \widetilde{\gamma}_s$ have total degrees at most d and logarithmic heights at most h.

Proposition 4 (Győry, E.)

If $\gamma_1, \ldots, \gamma_s$ are multiplicatively dependent, then there are integers $k_1, \ldots, k_s$, not all 0, such that

$$\gamma_1^{k_1} \cdots \gamma_s^{k_s} = 1, \quad \max_i |k_i| \le (d+2)^{\kappa^{r+s}} (h+1)^{s-1}$$

where κ is an effectively computable absolute constant > 1.

Unit equations vs. exponential equations

Theorem (Roquette, 1956)

Let A be a finitely generated domain over $\mathbb{Z}$. Then its unit group A^* is finitely generated, i.e., there is a finite set of generators $\gamma_1, \ldots, \gamma_s \in A^*$ such that $A^* = \{\gamma_1^{u_1} \cdots \gamma_s^{u_s} : u_i \in \mathbb{Z}\}$.

By Roquette's Theorem, the unit equation

$$(1) \qquad ax + by = c \quad \text{in } x, y \in A^*$$

can be rewritten as an exponential equation

$$(5) \qquad a\gamma_1^{u_1} \cdots \gamma_s^{u_s} + b\gamma_1^{v_1} \cdots \gamma_s^{v_s} = c \quad \text{in } u_1, \ldots, v_s \in \mathbb{Z}.$$

But as yet, no algorithm is known which for an arbitrary given finitely generated domain A over $\mathbb{Z}$ computes a finite set of generators for A^*.

So from an effective result on (5) one can not deduce an effective result on (1).

Idea of proof of Theorem 2

Let $A = \mathbb{Z}[z_1, \ldots, z_r]$. We can map

(1) $$ax + by = c \quad \text{in } x, y \in A^*$$

to S-unit equations in a number field by means of specializations

$$\varphi : A \to \overline{\mathbb{Q}} : z_i \mapsto \xi_i \in \overline{\mathbb{Q}} \quad (i = 1, \ldots, r).$$

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1. Apply 'many' specializations to (1) and apply the effective result of Győry-Yu to each of the resulting S-unit equations. This leads, for each solution x,y of (1) and each of the chosen specializations φ, to effective upper bounds for the logarithmic heights $h(\varphi(x))$ and $h(\varphi(y))$.

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2. View (1) as an equation over the algebraic function field $\mathbb{Q}(z_1, \ldots, z_r)$ and determine effective upper bounds for the function field heights $h_f(x)$, $h_f(y)$, using Stothers' and Mason's effective abc-Theorem for function fields.

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3. Combine the bounds from 1) and 2) with Aschenbrenner's theorem on linear equations over $\mathbb{Z}[X_1, \ldots, X_r]$, to get effective upper bounds for the sizes of representatives for x, y. $\square$

Work in progress (with Attila Bérczes, Kálmán Győry)

Effective results over finitely generated domains A (with effective upper bounds for the sizes of the solutions) for

- Thue equations $F(x, y) = \delta$ in $x, y \in A$ (F binary form in $A[X, Y]$, $\delta \in A \setminus \{0\}$);
- Hyper- and superelliptic equations $y^m = f(x)$ in $x, y \in A$ ($f \in A[X]$, $m \geq 2$);
- Discriminant form equations $\mathrm{Discr}_{L/K}(\alpha_1 x_1 + \cdots + \alpha_m x_m) = \delta$ in $x_1, \ldots, x_m \in A$ (K quotient field of A, L finite extension of K, $\alpha_1, \ldots, \alpha_m \in L$, $\delta \in A \setminus \{0\}$).

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